

The Evolving Threat of Internet Worms

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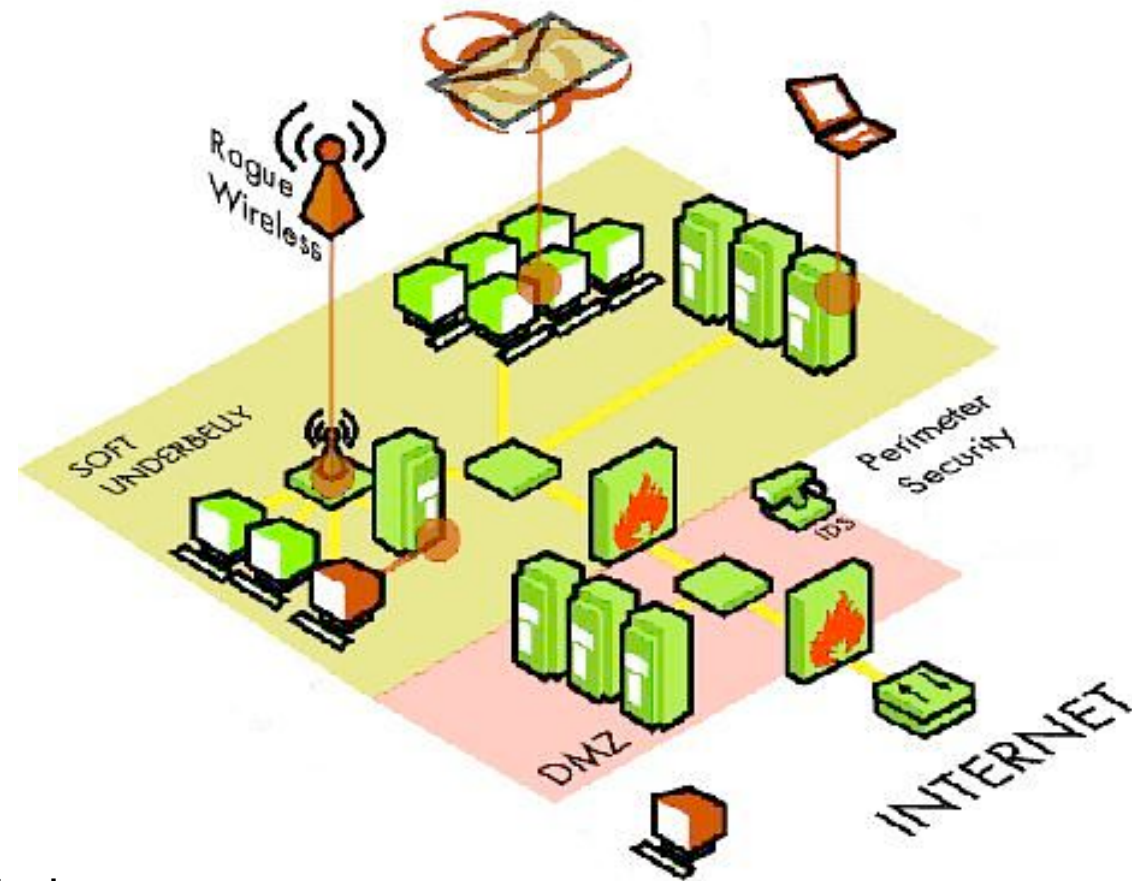


Why Worm Based Intrusions



- Relative ease
 - “Write once, run everywhere” promise can come true
- Penetration
 - Right past firewalls via laptops, find the weakest link
- Persistence
 - Worms keep working so you don’t have to
- Coverage
 - Attack everything eventually

Why Worms are Successful



- Missing patches
- Rogue access
- Missing access control

Evolution of Worm Threats



- Previously, worms were simple clones

- Worms have become more complicated systems
 - Multi-vector
 - Grow your potential target base
 - DDoS tool propagation
 - Utilize the army of machines
 - Dynamic
 - Thwart static detection mechanisms
 - Counterworms
 - Fight back with the same strategy

Multi-vector Worms



Goal is to thwart simple defenses and infect more machines

- Code Red vs. Nimda (2001)
 - Code Red: one attack vector (IIS)
 - Nimda: multiple attack vectors (IIS, mail, IE, open shares)
- Sircam (2001)
 - Mass mailer, also spread via open shares
- Blaster (2003)
 - MS-RPC or WebDAV attacks

DDoS Tool Propagation



Use the worm to attack an adversary

- Code Red (2001)
 - SYN flood against a static IP
- Blaster (2003)
 - SYN flood against a static domain
 - Variants carried a DDoS toolkit
- Sapphire, Welchia (2003)
 - The worm's spread is a DDoS

Dynamic Worm Appearances



Try and develop a worm with longevity by evading defenses

- Hybris (2000)
 - Used alt.comp.virus to spread code updates
- Lirva (2003)
 - Attempted to download new packages from website
- Sobig (2003)
 - Contacted website for next set of instructions

Counterworms



Fight the worm with a fast, scalable attack

- Code Green (2001)
 - Anti-Code Red worm
- Cheese (2001)
 - Anti-L1on worm
- Welchia (2003)
 - Anti-Blaster worm

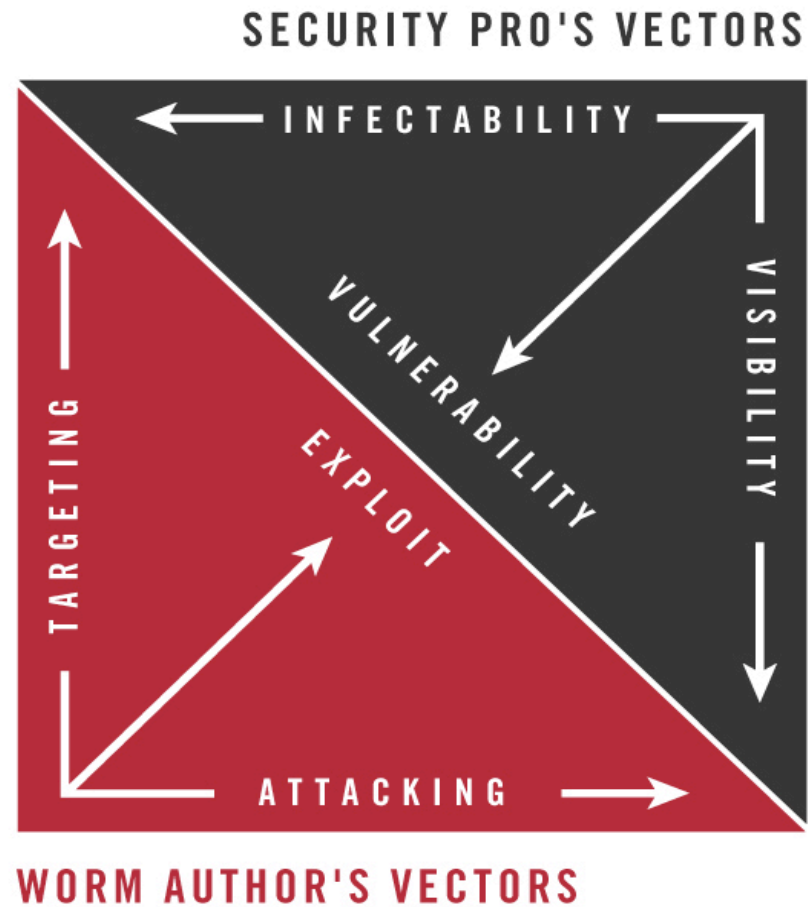
Cause more traffic and problems than they attempt to solve

Worm Authors Are Learning



- It's growing easier to build worms
 - Recycle an exploit, automation code, build, launch
- Use flexible targeting for DoS attacks
- No need to target multiple platforms
 - One platform works well enough
- Multiple infection vectors lead to longevity
 - Nimda still present two years later
- Local bias effective at enterprise penetration
 - Worms will be carried into the enterprise
 - Laptops, VPN connections

Vectors of Control



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Current Visibility Control



- Classic firewall strategy for the Internet
 - Minimally protect the DMZ
 - Maximally protect the internal network
- DMZ for exposed services
 - Control data flow between trusted, untrusted networks
- Hardened wall against internal, external networks

Classic Vulnerability Control



- ⊕ Minimized setups on system rollouts
 - Construct an image with minimal software
- ⊕ Patch maintenance
 - Worms typically attack known holes
- ⊕ Aggressive known vulnerability inventorying
 - Regular system inventories, comparisons against vulnerability databases (e.g. CVE)

Controlling Infectability



- Hardened systems
 - OS level changes
 - Non-executable stack
 - Permissions for any subsystem

- Hardened applications
 - Application configurations

- Strengthened configurations
 - Services and privileges for any system

Going Beyond the Firewall



- Traditional firewall configuration methods
 - Decide policy, install filters
 - Adjust by reading logs, tweak as needed
 - Broken applications or upset users

- Informed firewall configurations
 - Measure traffic, infer usage
 - Determine policy, install policy

- Assisted by Peakflow X

Intelligent Risk Assessment



- Traditional vulnerability scanners
 - Scan for a service, list machines offering that service
 - Banner grab, report service type, report potential vulnerabilities

- Usage, policy-aware vulnerability scanners
 - Scan for services, compare against usage and policy, report differences
 - Performed by Peakflow X

Combating Worms



- Minimize visibility
 - Tune access filters to a minimal set
 - Externally reachable
 - Internally used

- Minimize vulnerability
 - Track used services
 - Identify, remove unused services
 - Couple to strong patch management

Detecting Worms



⊕ Challenge

- In the face of dynamic behaviors, reliably detect the presence of a worm

⊕ Solution

- Every worm attempts to spread from host to host
- Specific forms of traffic will increase
- Not every host will have sent this traffic before
 - Example: web server becoming a web client

⊕ Therefore

- Detect the cascading change in host behaviors

Data Gathering for Worm Detection



- Blackhole networks
 - No background traffic
 - Collect attempts from worm trying random hosts

- Live enterprise networks
 - Traffic and relationship modeling

- Live backbone networks
 - Interface and topology statistics
 - Traffic modeling and analysis

Principles of Correlation Analysis



- Two types of correlations to qualify events
 - Auto-correlation
 - Frequency and sources for any single type of anomaly
 - Example: scan frequencies
 - Cross-correlation
 - Frequency and sources of related anomalies
 - Example: scans followed by traffic increases

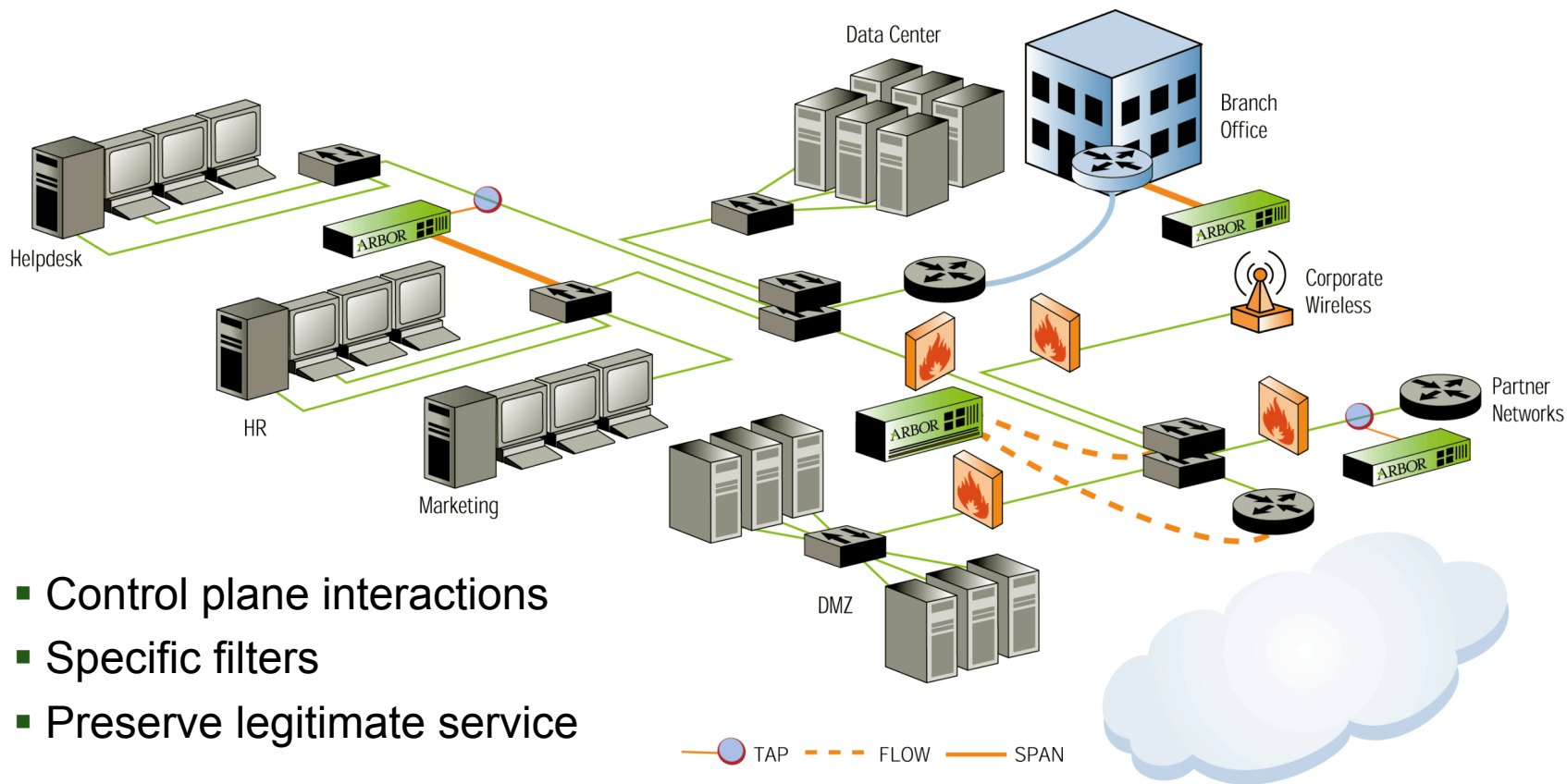
During worm outbreaks, these frequencies will increase from a growing number of hosts

Worm Detection by Peakflow X



- Uses correlation analysis
 - Partially based on an expert system
 - Extendable by the user via a filter language
- Produces a detailed report
 - Pattern of the worm's behavior
 - Hosts matching this pattern
 - Dynamically grows
 - Amount of traffic caused by the worm
- Couple to flow log for additional forensics

Safe Quarantine Interactions



- Control plane interactions
- Specific filters
- Preserve legitimate service

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Conclusions



- Worm authors are getting smarter
 - Worms are getting easier to write, more effective
- Worm detection mechanisms are getting more sophisticated and robust
- IDS and firewall mechanisms are advancing to develop worm defense techniques